



Course:CS603- Operating System

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| PROGRAMME: COMPUTERSCIENCE&ENGINEERING | DEGREE: B. TECH |
| COURSE: Operating Systems | SEMESTER: 6 CREDITS: 3 |
| COURSECODE: CS603 | COURSE TYPE: Theory |
| COURSE AREA/DOMAIN: Systems Software | CONTACTHOURS: 3 (weekly) |
| CORRESPONDINGLABCOURSE CODE (IFANY): CS693 | LABCOURSE NAME:Operating System Lab |

Course pre-requisites

| CODE | COURSE NAME | DESCRIPTION | SEM |
|-------|--|-------------------------------------|-----|
| CS201 | Basic Computation & Principles of Computer | Programming basics | II |
| CS303 | Computer Organisation | The components of a computer system | III |

Course Objectives

1. To develop an understanding of the functions underlying operating system
2. To learn the techniques for interfacing with the operating systems

Course Outcomes

1. Students would be able to explain the functions of an Operating System
2. Students would be able to understand the interaction between application software, Operating Systems and the hardware.
3. Students would be able to pick and choose the best OS design technique for a given problem
4. Student s would be able to design at least a module od OS.

ProgrammeOutcomes addressed in this course

- a. An ability to apply knowledge of mathematics, science, and engineering
- b. An ability to use the techniques, skills, and modern engineering tools necessary for engineering practice (k)



Syllabus

| UNIT | DETAILS | HOURS |
|------|--|-------|
| I | Introduction: Introduction to OS. Operating system functions, evaluation of O.S., Different types of O.S.: batch, multi-programmed, time-sharing, real-time, distributed, parallel. | 4 |
| II | System Structure: Computer system operation, I/O structure, storage structure, storage hierarchy, different types of protections, operating system structure (simple, layered, virtual machine), O/S services, system calls. | 3 |
| III | Processes : Concept of processes, process scheduling, operations on processes, co-operating processes, inter-process communication. | 3 |
| IV | Threads : Overview, benefits of threads, user and kernel threads. | 2 |
| V | CPU scheduling : Scheduling criteria, preemptive & non-preemptive scheduling, scheduling algorithms (FCFS, SJF, RR, priority), algorithm evaluation, multi-processor scheduling. | 3 |
| VI | Process Synchronization : Background, critical section problem, critical region, synchronization hardware, classical problems of synchronization, semaphores. deadlock prevention, deadlock avoidance, deadlock detection, recovery from deadlock. | 5 |
| VII | Deadlocks : System model, deadlock characterization, methods for handling deadlocks, deadlock prevention, deadlock avoidance, deadlock detection, recovery from deadlock. | 4 |
| VIII | Memory Management : Background, logical vs. physical address space, swapping, contiguous memory allocation, paging, segmentation, segmentation with paging. | 5 |
| IX | Virtual Memory : background, demand paging, performance, page replacement, page replacement algorithms (FCFS, LRU), allocation of frames, thrashing. | 3 |
| X | File Systems : file concept, access methods, directory structure, file system structure, allocation methods (contiguous, linked, indexed), free-space management (bit vector, linked list, grouping), directory implementation (linear list, hash table), efficiency & performance. | 4 |
| XI | I/O Management : I/O hardware, polling, interrupts, DMA, application I/O interface (block and character devices, network devices, clocks and timers, blocking and nonblocking I/O), kernel I/O subsystem (scheduling, buffering, caching, spooling and device reservation, error handling), performance. | 4 |
| XII | Disk Management : Disk structure, disk scheduling (FCFS, SSTF, SCAN,C-SCAN) , disk reliability, disk formatting, boot block, bad blocks. | 3 |



| | | |
|------|---|---|
| | Protection & Security : | |
| XIII | Goals of protection, domain of protection, security problem, authentication, one time password, program threats, system threats, threat monitoring, encryption. | 4 |



Gaps in the syllabus - to meet industry/profession requirements

| S.NO. | DESCRIPTION | PROPOSED ACTIONS | PO MAPPING |
|-------|---|------------------|------------|
| 1 | Batch processing History Explanation of how batch processing using punch cards used to happen in the early days. Need of code walkthrough to minimize compilation time and remove more errors | Extra Class | a. |
| 2 | How to Design a scheduler in different Operating system Understand how, we can design a scheduler | Extra Class | b |
| 3 | The role of hardware vis-à-vis software Understand which tasks are executed by the hardware and which by the software and the cost / time implication. | Extra Class | b |
| 4. | How to Design a IPC module Understand how, we can design a IPC Module | Extra Class | b |
| 5. | How to Design a Boot Strap Loader Understand how, we can design a Boot Strap Module | Extra Class | b |

Topics beyond syllabus/advanced topics

| S.NO. | DESCRIPTION | HOURS |
|-------|---|-------|
| 1 | Parallel and distributed processing Understanding the needs of a distributed system, specifically addressing distributed deadlock and consistent global snapshot. | 1 |

Web Source References

| S.NO. | URL |
|-------|---|
| 1 | http://web.stanford.edu/~ouster/cgi-bin/cs140-spring14/index.php |
| 2 | http://www.ics.uci.edu/~ics143/lectures.html |



Delivery/Instructional Methodologies

| S.NO. | DESCRIPTION |
|-------|--------------------------|
| 1 | Chalk and Talk |
| 2 | Study Material |
| 3 | Power Point Presentation |

Assessment Methodologies

| S.NO. | DESCRIPTION | TYPE |
|-------|------------------------|----------|
| 1 | Student Assignment | Direct |
| 2 | Tests | Direct |
| 3 | University Examination | Direct |
| 4 | Student Feedback | Indirect |



Course Plan

| S. NO. | Day | Module | Topic |
|--------|--------|--------|---|
| 1. | Day 1 | I | Introduction to OS. |
| 2. | Day 2 | | Operating system functions |
| 3. | Day 3 | | Evaluation of O.S. |
| 4. | Day 4 | | 4.Different types of O.S. Batch, Multiprogrammed |
| 5. | Day 5 | | Different types of O.S.: Time-sharing, real-time, distributed, parallel. |
| 6. | Day 6 | II | Computer system operation, I/O structure, storage structure, storage hierarchy. |
| 7. | Day 7 | | Different types of protections, operating system structure (simple, layered, virtual machine) |
| 8. | Day 8 | | O/S services, system calls |
| 9. | Day 9 | III | Concept of processes, process scheduling |
| 10. | Day 10 | | Operations on processes, Co-operating processes |
| 11. | Day 11 | | Inter-process communication. |
| 12. | Day 12 | IV | Overview, Benefits of threads, |
| 13. | Day 13 | | User and kernel threads. |
| 14. | Day 14 | V | Scheduling criteria, preemptive & non-preemptive scheduling, |
| 15. | Day 15 | | Scheduling algorithms (FCFS, SJF, RR, priority), |
| 16. | Day 16 | | Algorithm evaluation, multi-processor scheduling |
| 17. | Day 17 | VI | Background, critical section problem |
| 18. | Day 18 | | Semaphores and its applications |
| 19. | Day 19 | | Synchronization hardware |
| 20. | Day 20 | | Classical problems of synchronization |
| 21. | Day 21 | | Critical region, Monitor |
| 22. | Day 22 | VII | System model, Deadlock characterization |
| 23. | Day 23 | | Methods for handling deadlocks |
| 24. | Day 24 | | Deadlock prevention , Deadlock avoidance |
| 25. | Day 25 | | Deadlock detection, recovery from deadlock |
| 26. | Day 26 | VIII | Background, logical vs. physical address space |
| 27. | Day 27 | | Swapping, contiguous memory allocation |
| 28. | Day 28 | | Paging |
| 29. | Day 29 | | Segmentation |
| 30. | Day 30 | | Segmentation with paging |
| 31. | Day 31 | IX | Background, demand paging, performance, |
| 32. | Day 32 | | Page replacement, page replacement algorithms (FCFS, LRU), |
| 33. | Day 33 | | Allocation of frames, thrashing. |



| <i>S. NO.</i> | <i>Day</i> | <i>Module</i> | <i>Topic</i> |
|---------------|------------|---------------|--|
| 34 | Day34 | X | File concept, Access methods, Directory structure |
| 35 | Day 35 | | File system structure |
| 36 | Day 36 | | Allocation methods (contiguous, linked, indexed), free-space management (bit vector, linked list, grouping) |
| 37 | Day 37 | | Directory implementation (linear list, hash table), efficiency & performance |
| 38 | Day 38 | XI | I/O hardware, polling, interrupts, DMA |
| 39 | Day 39 | | Application I/O interface (block and character devices, |
| 40 | Day 40 | | Network devices, clocks and timers, blocking and nonblocking I/O), |
| 41 | Day 41 | | Kernel I/O subsystem (scheduling, buffering, caching, spooling and device reservation, error handling), performance. |
| 42 | Day 42 | XII | Disk structure, Disk scheduling (FCFS, SSTF, SCAN,C-SCAN) |
| 43 | Day 43 | | Disk reliability |
| 44 | Day 44 | | Disk formatting, Boot block, Bad blocks |
| 45 | Day 45 | XIII | Goals of protection, domain of protection, security problem |
| 46 | Day 46 | | Authentication, one time password |
| 47 | Day 47 | | Program threats, system threats |
| 48 | Day 48 | | Threat monitoring, encryption |



Assignment Set I

| S. NO. | Question | Assesses CO |
|--------|---|-------------|
| 1 | How contiguous allocation problem is solved using paging | CO.1 |
| 2 | Hardware need in paging | CO.2 |
| 3 | Why page size is always powers of 2? | CO.1 |
| 4 | Describe a mechanism by which one segment could belong to the address space of two different processes. | CO.2 |
| 5 | <p>In the IBM/370, memory protection is provided through the use of <i>keys</i>. A key is a 4-bit quantity. Each 2K block of memory has a key (the storage key) associated with it. The CPU also has a key (the protection key) associated with it. A store operation is allowed only if both keys are equal, or if either is zero. Which of the following memory-management schemes could be used successfully with this hardware?</p> <p>a. Bare machine b. Single-user system c. Multiprogramming with a fixed number of processes d. Multiprogramming with a variable number of processes e. Paging f. Segmentation</p> | CO.3 |
| 6 | <p>Consider a logical address space of eight pages of 1024 words each, mapped onto a physical memory of 32 frames.</p> <p>a. How many bits are there in the logical address? b. How many bits are there in the physical address?</p> | CO.2 |
| 7 | What is the effect of allowing two entries in a page table to point to the same page frame in memory? Explain how this effect could be used to decrease the amount of time needed to copy a large amount of memory from one place to another. What effect would updating some byte on the one page have on the other page? | CO.2 |
| 8 | Describe a mechanism by which one segment could belong to the address space of two different processes. | CO.2 |
| 9 | <p>Sharing segments among processes without requiring the same segment number is possible in a dynamically linked segmentation system.</p> <p>a. Define a system that allows static linking and sharing of segments without requiring that the segment numbers be the same. b. Describe a paging scheme that allows pages to be shared without requiring that the page numbers be the same.</p> | CO.3 |



10 Name two differences between logical and physical addresses. CO.3

Assignment Set II

| S. NO. | Question | Assesses CO |
|--------|--|-------------|
| 1 | Under what circumstances do page faults occur? Describe the actions taken by the operating system when a page fault occurs. | CO.1 |
| 2 | Assume that you have a page-reference string for a process with m frames (initially all empty). The page-reference string has length p ; n distinct page numbers occur in it. Answer these questions for any page replacement algorithms: a. What is a lower bound on the number of page faults? b. What is an upper bound on the number of page faults? | CO.1 |
| 3 | Which of the following programming techniques and structures are “good” for a demand-paged environment? Which are “not good”? Explain your answers. a. Stack b. Hashed symbol table c. Sequential search d. Binary search e. Pure code f. Vector operations g. Indirection | CO.2 |
| 4 | Consider the following page-replacement algorithms. Rank these algorithms on a five-point scale from “bad” to “perfect” according to their page-fault rate. Separate those algorithms that suffer from Belady’s anomaly from those that do not. a. LRU replacement b. FIFO replacement c. Optimal replacement d. Second-chance replacement | CO.2 |
| 5 | When virtual memory is implemented in a computing system, there are certain costs associated with the technique and certain benefits. List the costs and the benefits. Is it possible for the costs to exceed the benefits? If it is, what measures can be taken to ensure that this does not happen? | CO.3 |
| 6 | An operating system supports a paged virtual memory, using a central processor with a cycle time of 1 microsecond. It costs an additional 1 microsecond to access a page other than the current one. Pages have 1000 words, and the paging device is a drum that rotates at 3000 revolutions per minute and transfers 1 million words per second. The following | CO.3 |



statistical measurements were obtained from the system:

- 1 percent of all instructions executed accessed a page other than the current page.
- Of the instructions that accessed another page, 80 percent accessed a page already in memory.
- When a new page was required, the replaced page was modified 50 percent of the time.

Calculate the effective instruction time on this system, assuming that the system is running one process only and that the processor is idle during drum transfers.

- 7 Consider the two-dimensional array A: CO.2
- `int A[100][100];` where `A[0][0]` is at location 200, in a paged memory system with pages of size 200. A small process is in page 0 (locations 0 to 199) for manipulating the matrix; thus, every instruction fetch will be from page 0.
- For three page frames, how many page faults are generated by the following array-initialization loops, using LRU replacement, and assuming page frame 1 has the process in it, and the other two are initially empty?
- a. `for (int j = 0; j < 100; j++)`
 `for (int i = 0; i < 100; i++)`
 `A[i][j] = 0;`
- b. `for (int i = 0; i < 100; i++)`
 `for (int j = 0;`
- 8 Consider the following page reference string: CO.1
- 1, 2, 3, 4, 2, 1, 5, 6, 2, 1, 2, 3, 7, 6, 3, 2, 1, 2, 3, 6.
- How many page faults would occur for the following replacement algorithms, assuming one, two, three, four, five, six, or seven frames?
- Remember all frames are initially empty, so your first unique pages will all cost one fault each.
- LRU replacement
 - FIFO replacement
 - Optimal replacement
- 9 Suppose that you want to use a paging algorithm that requires a reference bit (such as second-chance replacement or working-set model), but the hardware does not provide one. Sketch how you could simulate a reference bit even if one were not provided by the hardware, or explain why it is not possible to do so. If it is possible, calculate what the cost would be. CO.1
- 10 You have devised a new page-replacement algorithm that you think may be optimal. In some contorted test cases, Belady's anomaly occurs. Is the new algorithm optimal? Explain your answer. CO.2



Assignment Set III

| S. NO. | Question | Assesses CO |
|--------|--|-------------|
| 1. | Design a CPU Scheduler which implements best scheduling algorithm. | CO.4,CO.3 |
| 2. | Design an IPC module for an Operating System. | CO. 4 |
| 3. | Design a Boot Strap Loader for an Operating System | CO.4 |